

Methodology for Evaluating Security Controls Based on Key Performance Indicators and Stakeholder Mission

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Abstract

Information security continues to evolve in response to disruptive changes with a persistent focus on information-centric controls and a healthy debate about balancing endpoint and network protection, with a goal of improved enterprise/business risk management. Economic uncertainty, intensively collaborative styles of work, virtualization, increased outsourcing and ongoing compliance pressures require careful consideration and adaptation. This paper proposes a Cyberspace Security Econometrics System (CSES) that provides a measure (i.e., a quantitative indication) of reliability, performance and/or safety of a system that accounts for the criticality of each requirement as a function of one or more stakeholders' interests in that requirement. For a given stakeholder, CSES accounts for the variance that may exist among the stakes one attaches to meeting each requirement. This paper introduces the basis, objectives and capabilities for the CSES including inputs/outputs as well as the structural and mathematical underpinnings.

1. Introduction

Good security metrics are required to make good decisions about how to design security countermeasures, to choose between alternative security architectures, and to improve security during operations. Therefore, in essence, measurement can be viewed as a decision aid. The lack of sound and practical security metrics is severely hampering progress in the development of secure systems.

Our Cyberspace Security Econometrics System (CSES) provides the following advantages over traditional measurement systems: (1) CSES accounts for the variances that exist among different stakeholders of the same system. Different stakeholders will typically attach different stakes to the

same requirement or service (e.g., a service may be provided by an information technology system or process control system, etc.). (2) For a given stakeholder, CSES accounts for the variance that may exist among the stakes one attaches to meeting each requirement. The same stakeholder may attach different stakes to satisfying different requirements within the overall system specification. (3) For a given compound specification (e.g., combination(s) of commercial off the shelf software and/or hardware), CSES accounts for the variance that may exist among the levels of verification and validation (i.e., certification) performed on the various components of the specification. In other words, the certification activity may produce higher levels of assurance across different components of the specification than others.

The relations between value based quantitative systems, risk assessment and other cyberspace research and applications areas is well founded [1-3]. These relations, characterized by CSES, are unavoidably involved with software and information system product and process technology, and their interaction with human values. CSES's rationale is strongly empirical, but includes new concepts in need of stronger theory. CSES uses risk considerations to balance information assurance discipline and flexibility, and to answer other key "how much is enough?" questions [4]. CSES will help to illuminate information technology policy decisions by identifying the quantitative and qualitative sources of cost and value associated with candidate decisions. In this paper we introduce the basis, objectives and capabilities for the CSES including inputs/outputs and the basic structural and mathematical underpinnings.

1.1. Qualities of Security Metrics

Some qualities of a good metric include: (1) ability to measure the right thing (e.g., supports the decisions that need to be made), (2) quantifiable (e.g., damages in dollars of profit loss), (3) capability to be measured precisely and accurately, (4) ability to be validated against ground truth, and (5) confidence level one has in the assertions made within the framework of the

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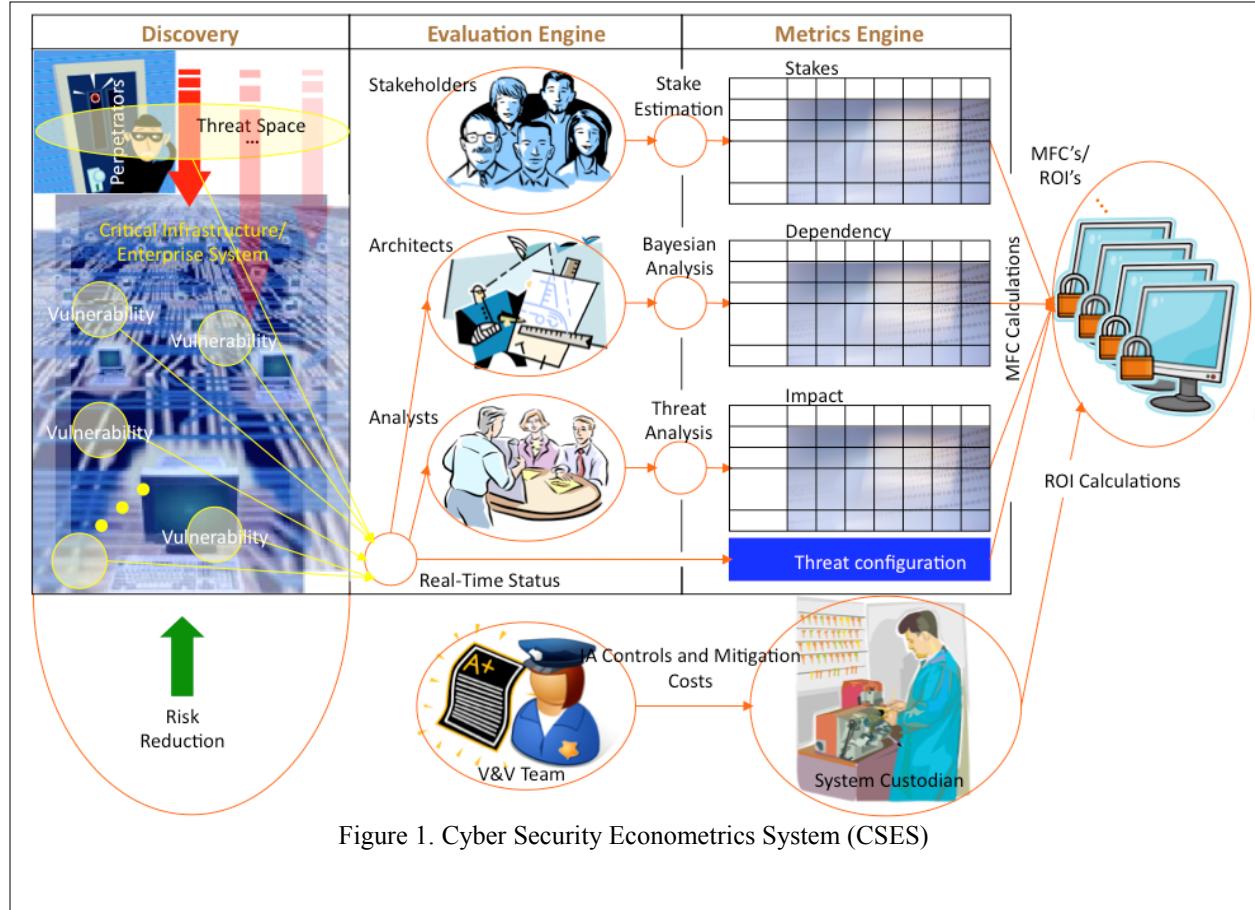
metric. To these criteria, one should add the following desirable properties: (1) inexpensive in time and cost to perform, (2) ability to be refereed independently, (3) repeatable so the outputs are independent of the analyst performing the measurement, and (4) scalable from small, single computer systems to large, nation-scale enterprise networks.

1.2. Quantifying Security Metrics

System security should be characterized, not by some abstract discrete scale, but rather by the very concrete determinant, mean failure cost (MFC) [5, 6]. MFC reveals how much each stakeholder stands to lose from mission value due to the lack of security. Subsequent use of this quantification enables us to derive an economic model that captures the tradeoffs involved in deploying security counter measures. Let us consider some fundamental pieces needed for evaluating security control.

1.3. Fundamental Concepts

Figure 1 shows essential input/output components and phases (i.e., discovery, evaluation, and metrics)



including data collection/analysis and consisting of the following:

System Stakeholders refers to any person or organization that has a stake in the operation of the system (i.e., users, operators of the system, hosts of the systems, etc.).

Security Specification used in the same way that correctness is a relative attribute (a system is correct with respect to its functional specification) and refers to a representation of the security attributes that a system must satisfy to be deemed secure.

Security Requirement used in the same way that a complex functional specification is typically composed of simpler components (representing elementary functional properties), and is composed of simpler security requirements [5, 7].

Mean Failure Cost used in the operational sense because the lack of security within the system may cause damage, in terms of lost productivity, lost business, lost data, resulting in security violations. We represent this loss by a random variable, and define MFC as the mean of this random variable [5]. As discussed further, this quantity is not intrinsic to the system, but varies by stakeholder.

2. Step-Wise Process of CSES

To estimate the mean failure cost of a system for a set of stakeholders, we need to initially identify and then maintain the following information: (1) the set of stakeholders of the system, and (2) the set of security specifications and thus security requirements that are to be required/provided by the system. (3) For each stakeholder and each security requirement, the stake that the selected stakeholder attaches to the selected service (or conversely, the cost that the stakeholder incurs if the service is disrupted). This information is provided by stakeholders. (4) For each component of a specific security requirement, the likelihood that the system provides that service as specified. This information is computed in light of the V&V measures (inspection, verification, testing, security measures, firewalls, vulnerability removal, threat mitigation, etc) that the system has undergone. In particular, estimating the likelihood of delivering a service requires that we analyze to what degree the components that are involved in delivering this service have been validated. The CSES process proceeds in three steps (Generation of Stakes Matrix, Dependency Matrix, and Threat Matrix), discussed in the next subsections [8, 9]. CSES encompasses not only failure costs but also mitigation costs, specifically verification costs [8].

2.1. Generation of Stakes Matrix

The MFC estimation depends on the following premises: (1) a stakeholder may have different stakes in different security requirements, and (2) a security requirement may carry different stakes for different stakeholders. The best way to represent this situation is through a two dimensional matrix, where the rows represent stakeholders, the columns represent security requirements and the entries represent stakes, as shown in Table 1.

The FC entry at row i , column j , represents the cost that stakeholder S_i would lose if the system failed to meet the security requirement R_j (i.e., also represented as $FC(S_i, R_j)$). The data in Table 1 is determined by

stakeholders. Each row is filled by the corresponding stakeholder, possibly in their own (possibly distinct) financial / economic terms (Dollars, Person Months, Euros, etc).

Using this computational infrastructure, we estimate the MFC of the system for a stakeholder i as the weighted sum of the stakes that he attaches to all

Table 1. *Stakes (ST) Matrix: Cost of failing a security requirement*

		Security Requirements				
		R1	R2	R3	...	Rn
Stakeholders	S1					
	S2					
	S3					
	...					FC_i^j
	Sm					
		Probabilities of Security Requirements Delivery				
						P^j

the services (security requirements) of interest, weighted by the probability of failing to deliver these services. Specifically,

$$MFC_i = \sum_{i=1}^n FC_i^j \times (1 - P^j).$$

We envision using this metric, not only off-line, to estimate the seaworthiness of the system for a particular stakeholder, but also for online monitoring of a real-time system. To the extent that the probabilities of service delivery (i.e., success of security requirement being met) may be profoundly altered by an ongoing attack or intrusion, this may adjust the calculations of the mean failure costs, and may push MFC for a particular stakeholder beyond

Table 2. Example of *Stakes (ST) Matrix showing requirements*

Stakeholder	Requirements	
	Requirement 1: Safety Requirement	Requirement 2: Timeliness
Passengers	Personal safety	Convenience, Scheduling
Airline Company	Liability for loss of life Reputation of airline	Reputation for timeliness / PR
Aircraft Manufacturer	Liability for loss of life Reputation of aircraft	Zero
Accident Insurance of Aircraft	Replacement Value of the Aircraft, Liability	Zero
Life Insurance of Passenger	Value of Life Insurance	Zero

acceptable values, triggering an alarm.

Recently, several techniques have emerged to characterize stakeholders and their respective interests. A variety of approaches extend from mapping security threats applying Computer Emergency Readiness Team (CERT) taxonomy; to stakeholders in the automotive industry [10, 11]; to investigations to secure vehicles against cyber attacks using the Defense-in-Depth paradigm [12]; to physically embedding information and communications technology (ICT) security into vehicles themselves [13]. Applications further extend to checklists [14] and improving the cyber incident mission impact in the military [15] and evaluating damage from cyber attacks [16]. The long-standing question, therefore, is, “How much security is enough?” To address this question, the subject domain of risk management must be addressed [4, 17]. When risk management is factored into the equation, classical approaches can be of benefit [18-20]. To illustrate the relevance of this **Stakes** matrix, and the idea of requirement-dependent failure cost, let us consider the example of a flight control system on board a commercial aircraft as shown in Table 2.

Stakeholders: The aircraft pilot; the passengers; the airline company; the aircraft manufacturer; the FAA; the insurance company that insures the aircraft; the insurance company that insures a passenger (life insurance); the EPA; etc.

Requirements: Adhering to Safety Requirements (maintaining the aircraft above stalling speed, ensuring never to reverse thrust in mid-air, ensuring landing gears are out before landing, etc); Adhering to FAA flight vector; Ensuring timely response to autopilot parameter adjustments; Maximizing fuel efficiency;

2), corresponding to two requirements: Safety and Timeliness.

Table 3. *Dependency* (DP) Matrix: Links requirements with components

		Components				
		C1	C2	C3	...	Ck
Requirements	R1					
	R2					
	R3					
	...					$\pi(R_i E_j)$
	Rn					

All these stakes are typically well known, and can be quantified financially. Using the **Stakes** Matrix, we estimate the MFC of Stakeholder *S* as follows:

$$MFC(S) = \sum_{R_i} P(R_i) \times FC(S, R_i).$$

Where $P(R_i)$ is the probability that the system fails to meet requirement R_i . In other words, the mean failure cost for stakeholder *S* is the sum, for all requirements, of the cost of failing these requirements weighted by the probability of failing them. We discussed above how to derive the Stakes table, that provide the terms $FC(S, R_i)$. In the next subsection, we discuss how to derive the probability terms.

2.2. Generation of Dependency Matrix

The question we address in this section (Table 3) is how to estimate the probability that a particular security requirement is violated in the course of operating the system for some period of time. The idea that we pursue here is to link the probability of failing a particular requirement with the probability of failure

Table 4. Example of *Dependency* (DP) Matrix: Links requirements with components

Components Requirements	Processing Component	Login Component	Secure Storage Component	User Profile Analysis
Freedom from Insider Threats	0.01	0.6	0.2	0.98
Protection of Critical Data	0.01	0.2	0.98	0.2
Access Control	0.01	0.98	0.4	0.1

Minimizing flight delay; Ensuring a smooth ride; Minimizing emission of greenhouse gases; etc.

Stakes: For the sake of illustration, we present below two sample columns of the Stakes table (Table

of a component of the system. The elucidation of this probabilistic link involves an analysis of the system’s architecture, to determine *which component contributes to meeting which requirement*. However, to

illustrate our method, we present a possible solution to this problem, under a simplifying hypothesis, which is that security violations affect no more than one component at a time.

We let $C1, C2, C3, \dots, Ck$ be the components of the system and we let Ei , for $1 \leq i \leq k$, be the event: *Failure of Component Ci*. Finally, we let E_{k+1} be the event: *No component has failed*. Then we can write (as an identity of probability calculus):

$$P(R) = \sum_{i=1}^{k+1} \pi(E_i) \times \pi(R | E_i).$$

Where the term $\pi(Ei)$ represents the probability of event Ei and the term $\pi(R|Ei)$ represents the probability of failing to satisfy requirement R , given hypothesis Ei (i.e., that event i has occurred). In general, it is fair to assume that in the absence of component failures, security requirements are vacuously satisfied, whence we write:

$$\pi(R | E_{k+1}) = 0.$$

As for the conditional probabilities of failing requirement R under the hypothesis of component failures, these may be estimated by the systems architect in light of the dependencies that exist between components and requirements. As an example, we consider the following situation:

- Security Requirement, R : Data item A must never be **altered** by an unauthorized user.
- The enterprise architecture and security policies provide that copies of data item A are stored in two components, $C1$, and $C2$.
- Then, we can infer:

$$\pi(R | E_1) = 0.$$

$$\pi(R | E_2) = 0.$$

Failure of any one component does not violate the requirement, since we can always retrieve the correct value from the unaffected component (*remember*: we assume that no more than one component at a time fails) and restore it.

We now consider a different situation:

- Security Requirement, R : Data item A must never be **accessed** by an unauthorized user.
- The architecture and security policies provide that copies of data item A are stored in two components, $C1$, and $C2$.
- Then, we can infer:

$$\pi(R | E_1) = 1.0.$$

$$\pi(R | E_2) = 1.0.$$

If one component fails in such a way that it can no longer control access to data item A, then Requirement R is violated.

We now consider a third hypothetical situation:

- Security Requirement, R : Access to the system must be contingent on proper user authentication through a login procedure.
- The architecture and security policies provide that two components, $C1$, and $C2$, can login an incoming user. Incoming users are dispatched at random to one or the other.
- Then, we can infer:

$$\pi(R | E_1) = 0.5.$$

$$\pi(R | E_2) = 0.5.$$

If one component fails then system access is secure only for those users dispatched to the other non-faulty component.

More generally, we assert that an analysis of the system architecture, by architecture subject matter experts, can lead to the derivation of conditional probabilities that link the probability of component failures with the probabilities of failing to meet specific requirements. This information can be represented in a two dimensional matrix, which we call the **Dependency** matrix.

Table 3 shows its structure, where the term $\pi(Ej)$ represents the probability of event Ej and the term $\pi(R|Ej)$ represents the probability of failing to satisfy requirement Ri , given hypothesis Ej (i.e., that event j has occurred). In Table 3, there exists a component event Ej for a requirement Ri where the probability of failure to satisfy requirement R exists ($\pi(Ri | Ej)$).

In the current airline example, the following would be an example of a generation of the **Dependency** matrix with respect to the passenger requirements of (1) safety and (2) timeliness with three distinct different components for each requirement. The results are illustrated in Table 4.

The probability of failing a requirement is obtained by the sum, for all components, of the conditional probabilities of failing that requirement, conditional on failure of the component, weighted by the probability of failure of the component. The conditional probabilities are given by the dependency matrix; as for the probabilities of component failures, as for the probabilities of component failures and are the subject of the next section.

2.3. Generation of Impact Matrix

The Impact Matrix shows the Component Failure versus Threats Relationship Grouping (Table 5). The probability of component failure depends on three factors: (1) the armor (e.g., Technical Controls or mitigations) that the component is provided with to protect against threats and to mitigate damage in the case of successful attacks. (2) Pattern of threats that the component is subjected to. In the same way that researchers of reliability define fault models that catalog faults under consideration, we must define threat models that catalog what threats or families of threats we want to protect against. An example of classification may include: insider threats; intrusions; denial of service threats; authentication threats; etc. (3) the degree of verification and validation that the component has undergone, be it through testing, inspection, static analysis, etc.

To assess the likelihood that a particular threat leads to the failure of a component, we consider a set of catalogued threats (or families of threats with common attributes), say $T_1, T_2, T_3, \dots, T_h$, and we consider the events $V_1, V_2, V_3, \dots, V_h, V_{h+1}$, where V_i , for $1 \leq i \leq h$, stands for: *Threat i has materialized*, and V_{h+1} stands for: *No threat i has materialized*. Because events V_i , for $1 \leq i \leq h+1$, are complementary (if we assume that no more than one threat materializes at a time), we can write:

$$\pi(E_i) = \sum_{j=1}^{h+1} \pi(V_j) \times \pi(E_i | V_j).$$

This equation links the probability of threat T_j (which is $\pi(V_j)$) to the probability of component failure for component C_i (which is $\pi(E_i)$). To apply this formula, we need to derive the conditional probabilities, which we propose to represent in a two dimensional matrix, that we call the ***Impact*** matrix in Table 5.

This ***Impact*** matrix is filled by component analysts and security experts, by assessing the impact that each

Table 5. *Impact* (IM) Matrix: Links threat relationship groupings

		Threats				
		T1	T2	T3	...	Th
Components	C1					
	C2					
	C3					
	...					$\pi(E_i V_j)$
	Ck					

type of threat may have on the operation of the component. Automated (e.g., in the airline example or in the case of cyber security the Common Vulnerability Scoring System [CVSS]) or manual (e.g., Subject Matter Experts [SMEs]) mechanisms will be utilized, as the quantification of the asset is critical for the methodology as depicted in Table 6.

2.4. Generation of Mitigation Costs Matrix -

Our quantitative model CSES encompasses not only failure costs, but also mitigation costs – specifically verification costs. Each requirement fulfilled or service delivered by the system depends on the correct operation of one or more system components. This dependency can be quantified by the statistical correlation between the failure of the

Table 6. Impact (IM) Matrix: Links components to threats

Threats Components	Insider Threats	Intrusions	Denial of Service Threats	Authentication Threats	No Threat
Processing Component	0.2	0.4	0.8	0.8	0.0
Login Component	0.2	0.2	0.2	0.2	0.0
Secure Storage Component	0.2	0.4	0.2	0.2	0.0
User Profile Analysis	0.2	0.1	0.1	0.1	0.0

component and the failure to deliver the service or fulfill the requirement. If we combine this dependency with the cost of verifying each component of the system, we can maintain an estimate of the probabilities of service delivery (discussed above) as a function of the effort invested in enhancing the dependability of the individual components. Maintaining this information can serve two purposes:

- First, to determine, at all times, which components must be enhanced first to improve overall stakeholder satisfaction.
- Second, to charge verification costs according to stakeholder benefit. For any particular verification measure, we charge stakeholders according to the gains they have achieved as a result of this measure (which are quantified by the reduction of their mean failure cost).

Table 7 illustrates the dependency between verification costs per component and verification cost per service, and links to the table above, which illustrates the dependency between the probabilities of service delivery (resulting from enhanced verification) and mean failure costs.

The verification cost by service can be estimated by the following formula:

$$VS_i = \sum_{i=1}^n D_i^j \times VC_j.$$

3. Estimating the Probability of Threats

If we review our discussions so far, we find many formulas of the form:

$$Y_i = \sum_{1 \leq j \leq m} A_i^j \times X^j, 1 \leq i \leq n,$$

where Y is a vector of size n , X is a vector of size m ,

Table 7. *Mitigation Costs (MC)* Matrix: Links requirements and components mitigation costs

		Components							
		C1	C2	C3	C4	C5			
Services	S1						Verification Cost by Service	VS1	
	S2							VS2	
	S3							VS3	
	S4							VS4	
	S5							VS5	
		Verification Cost by Component							
		VC1	VC2	VC3	VC4	VC5			

and A is a $n \times m$ matrix. We write this formula compactly as:

$$Y = A \circ X.$$

Specifically,

- The vector of mean failure costs (MFC , one entry per stakeholder) is given by the following equation:

$$MFC = ST \circ PR,$$

where ST is the **Stakes** matrix and PR is the vector of requirement failure probabilities (one entry per requirement),

- The vector of requirement failure probabilities is given by the following equation:

$$PR = DP \circ PE,$$

where DP is the **Dependency** matrix and PE is the vector of component failure probabilities (one entry per component),

- The vector of components failure probabilities is given by the following equation:

$$PE = IM \circ PV,$$

where IM is the **Impact** matrix and PV is the vector of threat emergence probabilities (one entry by type of threat).

By substitution, we find the equation that gives us vector of mean failure costs of all stakeholders as:

$$MFC = ST \circ DP \circ IM \circ PV.$$

Utilizing a user interface, the **Stakes** matrix (ST) is filled by stakeholders according to the stakes they have in satisfying individual requirements; the **Dependency** matrix (DP) is filled in by the system architect (i.e.,

cyber security operations and system administrators) according to how each component contributes to meet each requirement; the ***Impact*** matrix (*IM*) is filled by analysts according to how each component is affected by each threat.

The remaining question is how to fill the vector *PV* that represents the probability of emergence of the various threats that are under consideration? This is done empirically, by simulating and/or operating the system for some length of time and estimating the number of threats that have emerged during that time and continue to be refined as the system evolves. From these numbers, we infer the probability of emergence of all the threats during one hour of operation.

4. Application of CSES

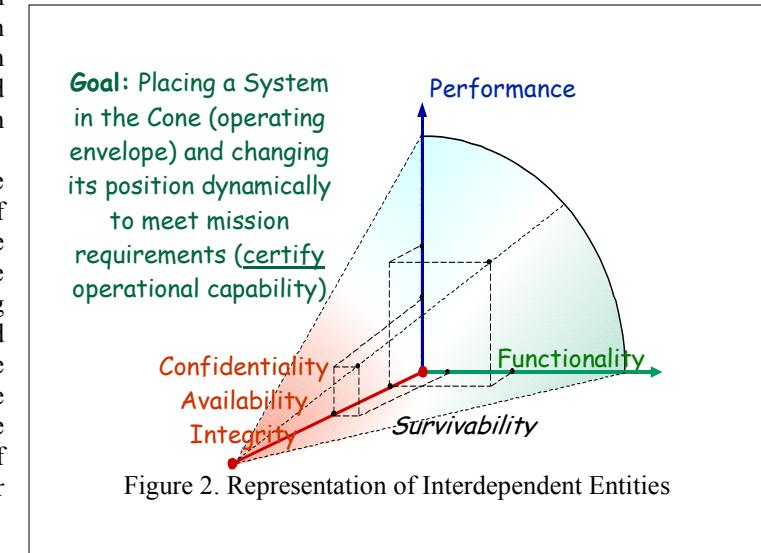
In this section, we provide a general set of guidelines for using the CSES.

4.1. Online Security Monitoring

In the *MFC* equation given above, the matrices *ST*, *DP*, and *IM* are static, and vector *PV* accounts for the analysts' estimate of threat emergence in the absence of any specific situational information. When this tool is coupled with sensors that detect anomalies or security violations, then it can dynamically re-compute the mean failure costs of stakeholders provided we know how new situational information impacts the vector *PV*. As an extreme example, if a threat of type *T1* has been confirmed, then vector *PV* is assigned the value $(1, 0, 0, 0, 0 \dots)$ and the new values of *MFC* are recomputed, giving stakeholders an estimate of their current mean failure cost if the system continues to operate. In this way, stakeholders may place the systems in the cone (i.e., protection/operating envelope) and changing its position dynamically, to meet mission requirements and to certify operational capability as shown in Figure 2 [21, 22].

4.2. Charging Verification & Mitigation Costs

Imagine that the custodian of the system wishes to submit a given component, say *CI*, to further verification and validation (V&V) or Information Assurance (IA) controls, and wishes to charge stakeholders for the V&V or modification of mitigation effort. As a matter of fairness, we wish to charge



stakeholders according to how much a better quality component impacts their bottom line. We argue that one way to do that is to estimate, for each stakeholder, the difference between the MFC under the current situation and the MFC after integration of mitigation technologies or V&V. How is this difference calculated? By considering how the modified IA controls or verification and validation of the component reduces the values of the failure probabilities in the row corresponding to that component in the impact matrix.

4.3. Integrating Quality Costs

From the standpoint of each stakeholder, the mean failure cost (which is the cost we expect to incur as a result of the lack of security) must be balanced against the cost of improving system security. Our mean failure cost model allows us to formulate the tradeoff of quality versus cost in terms of a return on investment equation. Specifically, a return on investment model is defined by the following parameters:

- An initial investment cost, say *IC*,
- An investment cycle (duration), say *T*,
- An return over the investment cycle, say *B(t)*, for $1 \leq t \leq T$, and
- A discount rate, say *d*.

Then the return on investment is given by the following formula:

$$ROI = -1 + \sum_{t=1}^T \frac{B(t)}{IC \times (1+d)^t}.$$

The formula identified in section 2.4 can be used to compute IC . We can estimate the benefit gained by stakeholder S during time period t by computing the difference between the mean failure cost with the current component and the mean failure cost (hopefully lower) with the validated component.

5. Summary

The CSES follows a defined process. The initial inputs (1) organization mission (and components thereof), (2) value of its objectives and assets if uninterrupted, and (3) the components of the enterprise system that support each mission component, are determined by stakeholders.

The customer, with assistance from SMEs, defines the criteria of a quantitative value of an asset. For example, the criteria may include:

- Financial basis (e.g., operational cost of downtime per unit of time defined with hardware/software costs, HVAC, staffing, etc versus profit); which is the quantitative measurement to be used within the CSES.
- Federal Information Security Management Act (FISMA) of 2002, customer derived value of assets per NIST 800-60, and/or FIPS 199/200 (February 2004, Standards for Security Categorization of Federal Information and Information Systems) dictated requirements.
- Customer defined requirements; acceptable and unacceptable impact levels against the value related to IA tenets of confidentiality, availability and integrity may also be examined in further program contract options to supplement the prototype.

The CSES process proceeds in three steps (Generation of Stakes Matrix, Dependency Matrix, and Threat Matrix). CSES encompasses not only failure costs but also mitigation costs, specifically verification costs. CSES provides:

- A framework for measuring the appropriate attributes that support the decisions necessary to (1) design security countermeasures, (2) choose between alternative security architectures, (3) respond to events such as intrusions or attacks and (4) improve security (including reliability and safety) during both design and operational phases.
- A comprehensive basis for choosing courses of action that have the highest risk reduction return on investment, i.e., reduce the most risks for the lowest cost.

The basis of CSES stems from and is consistent with the spirit of Value Based Software Engineering. CSES comprehends the different organizational

mission needs for all stakeholders, including reliability and safety. CSES identifies information assurance controls and mitigation costs as an investment toward assuring mission success.

6. Future Work

Our future plans will draw from our experience to provide a good example of a system with credible and distinct stakeholders, a rich security requirements structure (i.e., set of), a non-trivial architecture with intertwined security components and functional (processing) components as well as a rich catalog of families of threats.

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